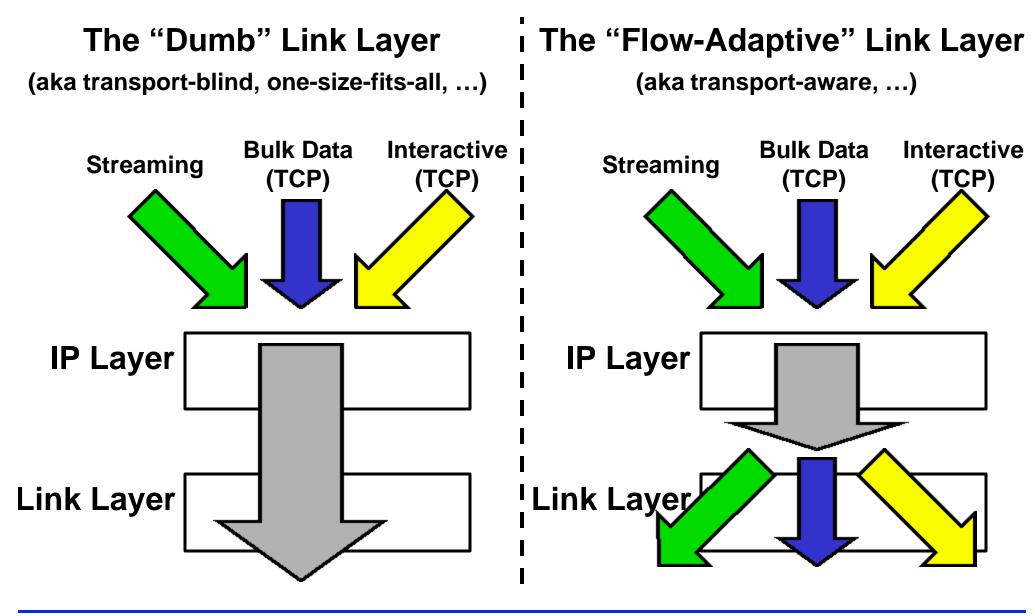
# Can Wireless Preserve the E2E Argument?

Dumb vs. Flow-Adaptive Link Layers (LL) Low vs. High LL ARQ Persistency for TCP

**Reiner Ludwig** 

Ericsson Research

# Link Layer Design Philosophies

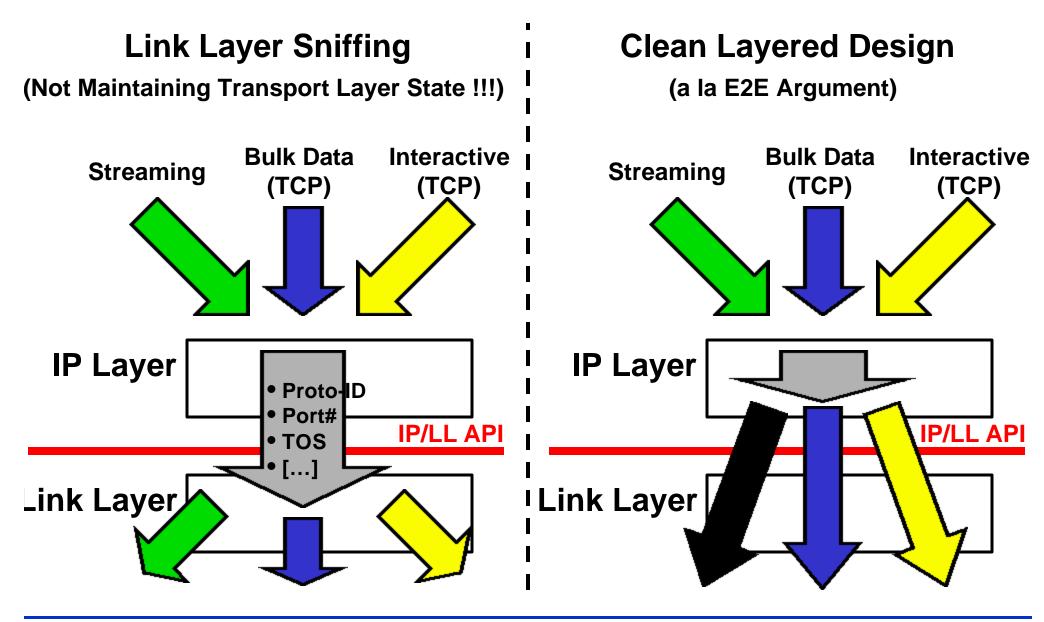


## Wireless Link Layers SHOULD be Flow-Adaptive

- Flow-Adaptive Makes Little Sense for Wireline
  - ⇒ Because: Wireline Link Layers have No Knobs for Tuning (not needed!)

- Flow-Adaptive Makes Lots of Sense for <u>Wireless</u>
  - ⇒ Because: Wireless Link Layers have Many Knobs for Tuning: FEC, Interleaving, ARQ, Power Control, ...
  - ⇒ Allows to Adapt Knobs to Flow's QoS Requirements
  - ⇒ Spectrum Efficiency
  - ⇒ Power (Battery) Efficiency

# How to Implement a Flow-Adaptive Link Layer



# BUT: E2E Argument Promotes Dumb LL ...

- "Everything should be true ind-points. The network included to the second should remain dumb."
  - ⇒ J. H. Saltzer, D. P. Reed, D. D. Clark, "End-To-End Arguments in System Design", ACM Transactions on Computer Systems, Vol. 2, No. 4, November 1984.

#### • E2E Argument:

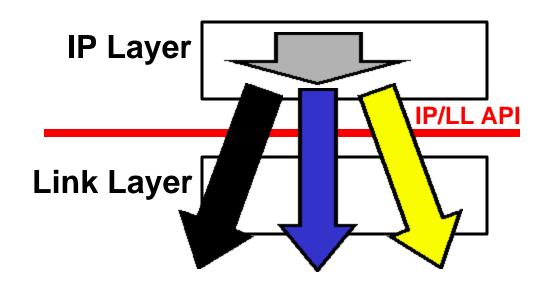
"[Link layer error control is] an incomplete version of the function provided by the communication system [that] may be useful as a performance enhancement".

# BUT: LL Sniffing is Layer Violation ...

True! On the other hand ...

- 1. Trade-off: Pragmatism/Performance vs. "Beauty of Design"
  - ⇒ If LL Sniffing (Layer Violation) was such a Concern: "Call the Layer-Police to Put the ROHCers into Jail" :-)

- 2. We have an alternative:
  - ⇒ Extended IP/LL API New PILC Work Item?



# BUT: Flow-Adaptive Breaks with IPsec ...

Partly True ...

- 1. People that are so Paranoid to use IPsec Gladly Trade Performance for Security.
  - ⇒ People who are less Paranoid Should Use TLS.
- 2. DS-field is unencrypted
- 3. IPsec-friendly Solution Possible (unencrypted TOS IP-Option?)

# Link Layer ARQ Persistency for Reliable Flows (TCP)

- Assume Flow-Adaptive LL, i.e, TCP flows are separated
- Assume LL ARQ is Possible
  - ⇒ Not the case on uni-directional links (e.g., some satellite links)
- LL ARQ Persistency for TCP?
- Definition of "LL ARQ Persistency":

The Time (in milliseconds) the LL Delays a Single IP Packet in an Attempt to Successfully Transmit it Across the Link.

#### BUT: We do Not Need LL ARQ ...

- - ⇒ Optimal Frame Size on some Wireless Links is less than 100 bytes (e.g., GSM, IS95, GPRS, UMTS)
  - ⇒ IPv6's Minimum MTU is 1280 Bytes!
  - ⇒ Might Work for Satellite Links: Optimal Frame Size >> 1280 Bytes

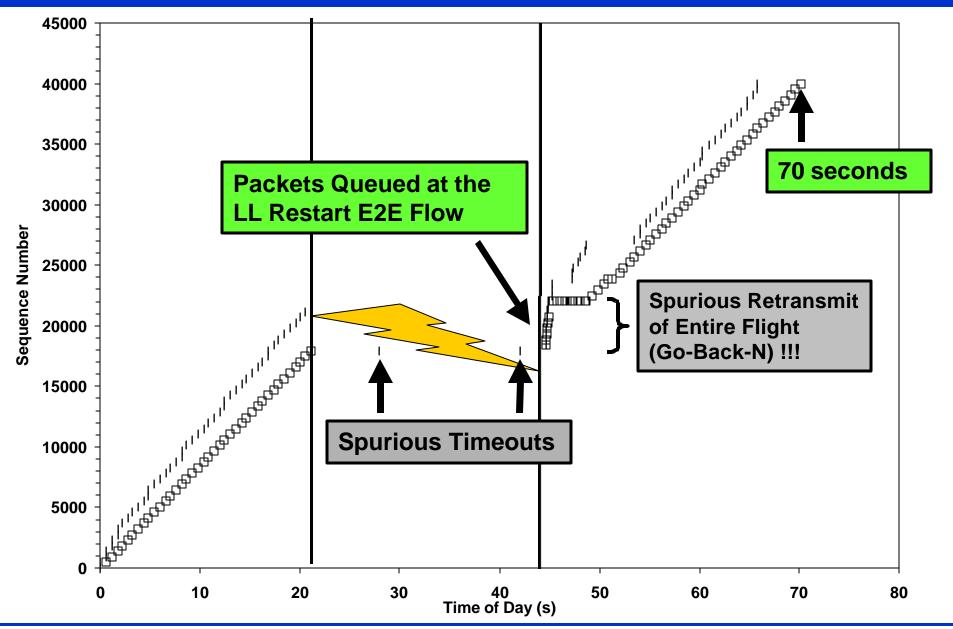
## Use Highly Persistent LL ARQ for TCP

- More Precisely, LL ARQ SHOULD try for up to 64 seconds (TCP's MAX-RTO) to Transmit a TCP Packet!
- This is NOT Saying: Unbounded Queues!
  - ⇒ Queues Need to Remain Small (Active Queue Management)
  - ⇒ If Queue Beyond Threshold **▶** "Drop From Front"
  - ⇒ Early Congestion Signal
- This is NOT Saying: Hop-By-Hop Instead of E2E Reliability!
  - ⇒ E2E Argument:
    - "[Link layer error control is] an incomplete version of the function provided by the communication system [that] may be useful as a performance enhancement".

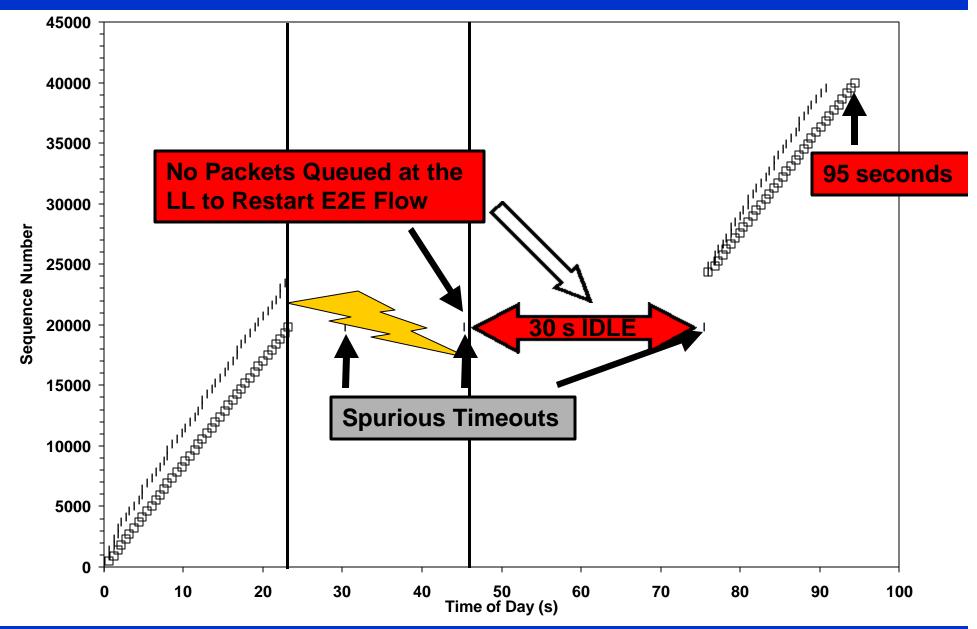
# Why Such a High LL ARQ Persistency?

- First of all, High Delays Due to LL ARQ are Rare
  - ⇒ Typically < 1 second (excluding transmission delay)
  - ⇒ Mainly Occurs During Transient Link Outages
- Most Spectrum & Energy (Battery) Efficient
  - ⇒ If the LL Can't Do it, TCP can't Do it!
  - ⇒ Discarding a Packet that Already Made it 90% Across the Link Makes No Sense!
  - ⇒ Measurements over GSM with LL ARQ Disabled and an MTU of 1500 Bytes Show up to 18% Undelivered Packets (Discarded by PPP due to CRC Error)
  - ⇒ RFC2914: "Congestion Collapse Due to Undelivered Packets"
- Robustness Against Link Outages
  - ⇒ No Need for an "ICMP-Link-Outage Agent" at the Basestation

# Link Outage & High LL ARQ Persistency



# Link Outage & Low LL ARQ Persistency



# **BUT: Spurious Timeouts ...**

True, they Force TCP into Go-Back-N. On the other hand ...

- 1. Likely to be Solved in TSV WG
  - ⇒ Eifel Algorithm
- 2. Go-Back-N Often Less Harmful than Waiting for Long RTO
  - ⇒ See Last 2 Slides

#### **BUT: Inflated RTO ...**

True! On the other hand ...

- 1.RTO Decays Quickly after an RTT Spike; especially when Timing Every Packet (Timestamp Option).
- 2. If the Path's RTT Varies Largely, RTO should be Inflated, i.e., should be conservative.

## BUT: Head of Line Blocking ...

Not True, as long as ...

... we Allow the LL to Perform Out-Of-Order Delivery Between Flows.

- ⇒ Requires LL Per-Flow Operations (Not Per-Flow State!)
- ⇒ However, No Scaling Concern on Last/First-Hop Links!

#### **A Word on TCP Proxies**

• TCP-Throughput = 
$$\frac{1}{RTT} \times \frac{1}{\sqrt{p}} \times C$$

- Flow-Adaptive LL + Highly Persistent LL ARQ for TCP
  - ⇒ Eliminates Non-Corgestion Packet Losses on Wireless Link
  - $\Rightarrow$  No Need for a Proxy to Avoid Influence on p
  - ⇒ For High Latency Links, a Proxy Might be Needed to Avoid Influence on *RTT*

#### A Word on Robust TCP/IP Header Compression

- Flow-Adaptive LL + Highly Persistent LL ARQ for TCP
  - ⇒ Eliminates Non-Congestion Packet Losses on Wireless Link
  - ⇒ No Losses Between Compressor & Decompressor
  - ⇒ No Need for Robustness in TCP/IP Header Compression Scheme!
  - ⇒ Only Things Left to do for ROHC WG: Compression of SYNs, FINs & TCP Option Fields (Timestamp, SACK, …)

# The Message

1. Wireless Link Layers SHOULD be Flow-Adaptive

2. Highly Persistent LL ARQ for TCP (all fully-reliable flows)

3. If 1. not feasible, e.g., due to IPsec, Low Persistent LL ARQ (< 100 ms ?) SHOULD be Operated for All Flows

#### Conclusion

# Can Wireless Preserve the E2E Argument?



# The E2E Argument is (Still) THE Guideline Leading to Well Designed Wireless Link Layers